

Advanced Automata Theory

Exercise Sheet 7

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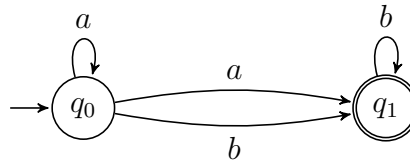
TU Braunschweig
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Out: May 24

Due: May 29, 12:00

Exercise 1: NBA Complementation

Consider the NBA A over $\Sigma = \{a, b\}$ below:



Use Büchi's complementation method discussed in class to compute $L(A)$ and $\overline{L(A)}$.

Exercise 2: Equivalence

Consider an NBA A , two classes $[u]_{\sim_A}$ and $[v]_{\sim_A}$ of \sim_A , and $w \in [u]_{\sim_A} \cdot [v]_{\sim_A}$ an ω -word. Show that if $w \in L(A)$ then $[u]_{\sim_A} \cdot [v]_{\sim_A} \subseteq L(A)$.

Exercise 3: Muller Automata

A Nondeterministic Muller Automaton (NMA) is a tuple $A = (Q, \Sigma, \delta, q_0, F)$. The first four components are as in Büchi automata. $F = \{Q_F^1, \dots, Q_F^n\} \subseteq \mathcal{P}(Q)$ is a set of sets of states instead of a single set of states. The idea is to accept a run if the set of states that occur infinitely often matches one of the Q_F^i exactly. Formally, a run r of A is accepting if $\text{Inf}(r) \in F$ where $\text{Inf}(r)$ is the set of states that are visited infinitely often in r . As for Büchi automata, we call A a Deterministic Muller Automaton (DMA) if for each $q \in Q$ and $a \in \Sigma$ there is exactly one state $q' \in Q$ such that $(q, a, q') \in \delta$.

- Given an NBA A , show that there is an NMA A_{NMA} such that $L(A_{NMA}) = L(A)$.
- Show that DMA are strictly more expressive than DBA.
- Given a DMA A , show that there is an NBA A_{NBA} such that $L(A_{NBA}) = L(A)$.
- Prove that DMA are closed under complement, i.e. for every DMA A there exists a DMA \bar{A} with $L(\bar{A}) = \overline{L(A)}$.