

Exercise Sheet 8

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Due: Tue, Jan 06

Exercise 8.1 Verifying Operating Systems

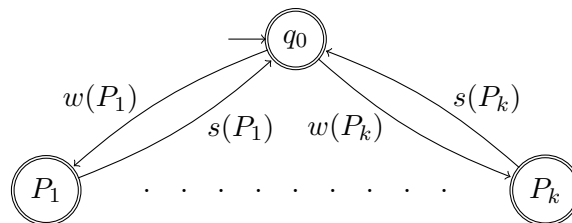
Our goal is to verify an operating system that runs k processes and has a scheduler. Consider the following Büchi automata:

- $A_{OS} := A_{P_1} \parallel \dots \parallel A_{P_k}$: Describes the behaviour of the operating system, where A_{P_i} represents the behavior of process P_i .
- A_{Sched} : Describes the scheduling strategy.
- A_{Prop} : Describes a property to be checked.

Our verification task amounts to solving the following model checking problem:

$$L(A_{OS}) \cap L(A_{Sched}) \subseteq L(A_{Prop}).$$

To solve this problem in a general way, we introduce a *most general* scheduling Büchi automaton A_{MG} that allows for arbitrary behaviours of the scheduler:



The scheduler can randomly *wake up* ($w(P_i)$) and *suspend* ($s(P_i)$) processes and the processes only work when awake. Unfortunately, this general scheduler is not *fair*: it does not necessarily wake up each process infinitely often.

- (a) Modify A_{MG} to a *fair* automaton A_{MGF} that wakes up every process infinitely often. Keep A_{MGF} as general as possible, do not implement a concrete scheduling strategy.
- (b) Present an automaton A_{RR} that describes the *Round Robin* scheduling strategy. What is the relationship between $L(A_{RR})$ and $L(A_{MG})$ respectively $L(A_{MGF})$?
- (c) Why can you conclude $L(A_{OS}) \cap L(A_{RR}) \subseteq L(A_{Prop})$ from $L(A_{OS}) \cap L(A_{MGF}) \subseteq L(A_{Prop})$?

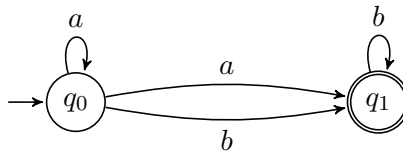
Exercise 8.2 NBA Emptiness and Membership

Let A be an NBA and wv^ω be an ω -word. Give algorithms that decide whether:

$$L(A) = \emptyset \qquad wv^\omega \in L(A).$$

Exercise 8.3 NBA Complementation

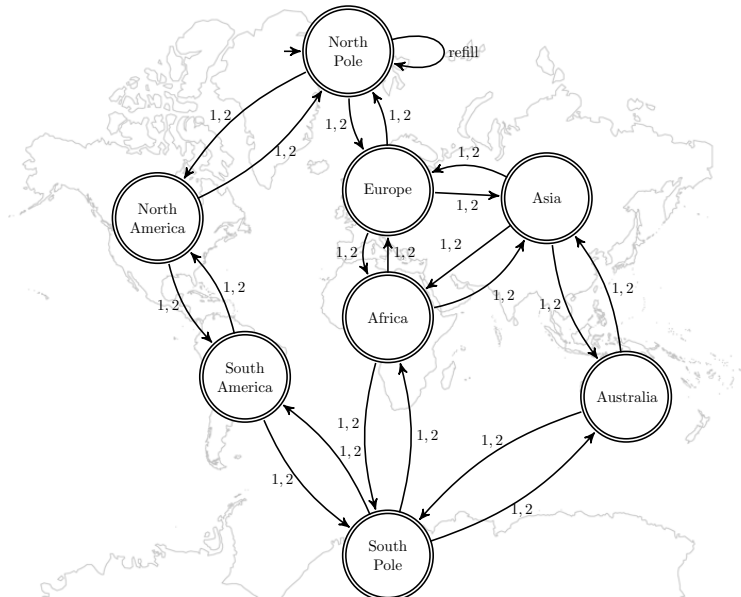
Consider the NBA A over $\Sigma = \{a, b\}$ below:



Use Büchi's complementation method discussed in class to compute $L(A)$ and $\overline{L(A)}$.

Exercise 8.4 Travelling Santa

Santa has decided to swap his traditional sleigh for the brand new *Christmas Racer 3000*, equipped with the state of the art automata driven navigation system *Rudolph Go v0.99beta*. The automaton $A_{Rudolph}$ controlling *Rudolph* is depicted below:



Each time Santa lifts off, *Rudolph* randomly chooses a neighbouring continent. Santa drops one or two presents every time he enters a region and he can refill at the north pole. Currently, *Rudolph* cannot prevent Santa from running out of presents. Please help Santa by upgrading *Rudolph* to version 1.0:

- Determine the minimum present capacity of the sleigh needed not to run empty.
- Give an NBA A_{Sleigh} modelling how the number of presents in Santa's sleigh changes.
- Explain how one can use $A_{Rudolph}$ and A_{Sleigh} to create a controller that guarantees Santa to never run out of presents.
- How can you modify the controller so that all continents are visited infinitely often?