

Concurrency theory

Exercise sheet 1

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Due: November 5

Submit your solutions until Tuesday, November 5, during the lecture. You may submit in groups up to three persons.

Exercise 1: The Ackermann function

a) The three-argument Ackermann function φ is defined recursively as follows.

$$\begin{aligned}\varphi: \mathbb{N}^3 &\rightarrow \mathbb{N} \\ \varphi(m, n, 0) &= m + n \\ \varphi(m, 0, 1) &= 0 \\ \varphi(m, 0, 2) &= 1 \\ \varphi(m, 0, x) &= m \quad \text{for } x > 2 \\ \varphi(m, n, x) &= \varphi(m, \varphi(m, n - 1, x), x - 1) \quad \text{for } n > 0 \text{ and } x > 0\end{aligned}$$

Formally prove the following equalities (e.g. using induction):

$$\varphi(m, n, 0) = m + n, \quad \varphi(m, n, 1) = m \cdot n, \quad \varphi(m, n, 2) = m^n.$$

b) Nowadays, one usually considers the following two-parameter variant.

$$\begin{aligned}A: \mathbb{N}^2 &\rightarrow \mathbb{N} \\ A(0, n) &= n + 1 \\ A(m, 0) &= A(m - 1, 1) \quad \text{for } m > 0 \\ A(m, n) &= A(m - 1, A(m, n - 1)) \quad \text{for } m > 0 \text{ and } n > 0\end{aligned}$$

For example, we have

$$A(1, 2) = A(0, A(1, 1)) = A(0, A(0, A(1, 0))) = A(0, A(0, A(0, 1))) = A(0, A(0, 2)) = A(0, 3) = 4.$$

Similar to this computation, write down a full evaluation of $A(2, 3)$.

Exercise 2: Konigs Lemma

A finitely branching tree \mathcal{T} is simply defined as a connected directed graph which satisfies the following properties.

- Every vertex in the tree is either a root node, internal node or a leaf node.
- There is exactly one root node, which has no incoming edges and has finitely many outgoing edges.
- Every internal node has exactly one incoming edge and finitely many outgoing edges.
- Every leaf node has exactly one incoming edge and no outgoing edges.

Prove that every finitely branching infinite tree has an infinite path. A path is a sequence of vertices in which the adjacent vertices are related by an edge.

Exercise 3: Petri net constructions

- a) Let (N, M_0, M_f) be a Petri net. Explain how to construct a Petri net (N', M'_0, M'_f) with $M'_0(p) = 0$ for all places but a single place p' with $M'_0(p') = 1$ and $M'_f(p) = 0$ for all places such that $M_f \in R(N, M_0)$ iff $M'_f \in R(N', M'_0)$. Recall that $R(N, M_0)$ denotes the set of all reachable markings of N , starting from M_0 .
- b) Let (N, M_0, M_f) be a Petri net. Explain how to construct a Petri net (N', M'_0, M'_f) such that M_f is coverable from M_0 in N iff M'_f is reachable from M'_0 in N' .
- c) Construct a Petri net N with only 3 places, a marking M_0 and markings $M_{c \wedge r}$, $M_{\neg c \wedge \neg r}$ and $M_{c \wedge \neg r}$ such that
- $M_{c \wedge r}$ is reachable and coverable from M_0 ,
 - $M_{\neg c \wedge \neg r}$ is neither reachable nor coverable, and
 - $M_{c \wedge \neg r}$ is coverable, but not reachable.

In each part of this exercise, argue briefly that your construction is correct.