Games with perfect information		
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Prof. Dr. Roland Meyer		Summer term 2017
Out: April 26		Due: May 2

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You can submit your solutions on Tuesday, May 2, at the beginning of the exercise classes (since Monday, May 1, is a public holiday.) I will grade and return the submissions as soon as possible!

Assume that we are in the same setting as in Section 6, Büchi and CoBüchi games: $G = (V_{\Box} \cup V_{\odot}, R)$ is a finite, deadlock-free game arena, and $B \subseteq V$ is a winning set. Recall the definitions of the sets B^i and P^i :

> $B^0 = B$, $P^{i} = V \setminus \operatorname{Attr}_{\bigcirc}(B^{i}),$ $B^{i+1} = B \setminus \operatorname{CPre}_{\square}(P^{i}).$

Recall that the controlled predecessors of a set $X \subseteq V$ were defined as

$$\mathsf{CPre}_{\stackrel{\wedge}{\succ}}(X) = \left\{ x \in V_{\stackrel{\wedge}{\succ}} \mid \exists (x,y) \in R \colon y \in X \right\} \cup \left\{ x \in V_{\stackrel{\wedge}{\succ}} \mid \forall (x,y) \in R \colon y \in X \right\}$$

and that using this definition, the attractor can be defined as follows.

$$\begin{aligned} &\operatorname{Attr}_{\stackrel{i}{\rightarrowtail}}^{0}(X) = X \\ &\operatorname{Attr}_{\stackrel{i}{\leadsto}}^{i+1}(X) = \operatorname{Attr}_{\stackrel{i}{\leadsto}}^{i}(X) \cup \operatorname{CPre}_{\stackrel{i}{\leadsto}}\left(\operatorname{Attr}_{\stackrel{i}{\leadsto}}^{i}(X)\right) \end{aligned}$$

Exercise 1: Proof of Lemma 6.9

Prove Lemma 6.9 from the lecture notes:

The sets B^i form a descending, the sets W_i form an ascending chain. There is an index $m \in \mathbb{N}$ such that the chains simultaneously become stationary.

$$B = B^{0} \supseteq B^{1} \supseteq \dots \supseteq B^{m} = B^{m+1} = \bigcap_{i \in \mathbb{N}} B^{i}$$
$$P^{0} \subseteq P^{1} \subseteq \dots \subseteq P^{m} = P^{m+1} = \bigcup_{i \in \mathbb{N}} P^{i}$$

Exercise 2: A more intuitive definition of recurrence sets

In this exercise, we give a more intuitive definition of the recurrence sets B^i , and we prove that it is equivalent to the definition in the lecture.

We need a slightly modified attractor construction:

$$A^{0}_{\stackrel{}{\not\sim}}(X) = \emptyset$$
$$A^{i+1}_{\stackrel{}{\not\sim}}(X) = A^{i}_{\stackrel{}{\not\sim}}(X) \cup \mathsf{CPre}_{\stackrel{}{\not\sim}}(A^{i}_{\stackrel{}{\not\sim}}(X) \cup X)$$
$$\mathsf{Attr}^{+}_{\stackrel{}{\not\sim}}(X) = \bigcup_{i \in \mathbb{N}} A^{i}_{\stackrel{}{\not\sim}}(X)$$

Now we give an alternative definition for the sets B^{i} , here called B_{ex}^{i} :

$$B_{ex}^{0} = B$$
$$B_{ex}^{i+1} = B \cap \operatorname{Attr}_{\bigcirc}^{+} \left(B_{ex}^{i} \right)$$

- a) Describe the difference between $Attr^+_{\preceq}(B)$ and $Attr_{\preceq}(B)$ in your own words. (Not more than 2 sentences please.)
- b) Formally prove using induction on *i* that $A^{i}_{\preceq}(B) \cup B = \operatorname{Attr}^{i}_{\preceq}(B)$ for all $i \in \mathbb{N}$ and conclude $\operatorname{Attr}^{+}_{\preceq}(B) \cup B = \operatorname{Attr}_{\preceq}(B)$.
- c) Formally prove using induction on *i* that $B^i = B^i_{ex}$ for all $i \in \mathbb{N}$.

Hint: In the induction step, you essentially need to prove

$$V \setminus \operatorname{Attr}_{\bigcirc}^{+}(B^{i}) = \operatorname{CPre}_{\square}\left(V \setminus \operatorname{Attr}_{\bigcirc}(B^{i})\right).$$

Part b) of this exercise is crucial for proving this statement.