

16. Fixed Points

- Goal:
- Introduce basic notions from lattice theory
 - Present the main fixed point theorems.

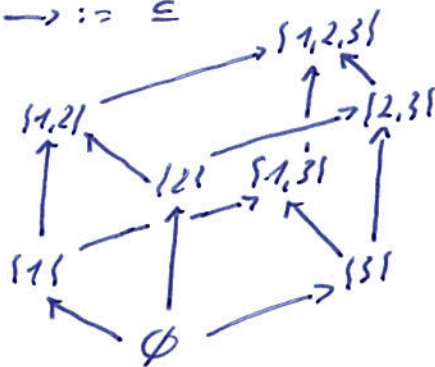
16.1 Complete Lattices

- Goal:
- Define partial orders, join and meet, lattices, complete lattices, bottom and top.

- Observation:
- (\mathbb{N}, \leq) is totally ordered, every two elements are comparable.
 - Some domains are only partially ordered:

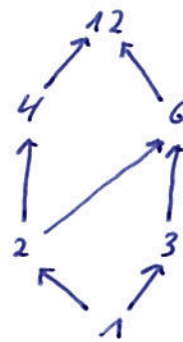
Subsets of $\{1, 2, 3\}$,

$\rightarrow := \subseteq$



$\{1, 2\}$ and $\{2, 3\}$
are incomparable

Divisors of 12
 $\rightarrow := |$



2 and 3 are
incomparable.

Definition:

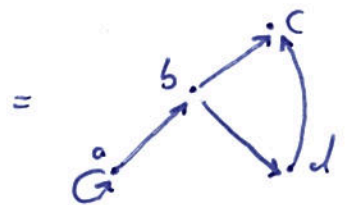
A partial order (D, \leq) consists of a set $D \neq \emptyset$

and a reflexive, transitive, and anti-symmetric relation $\leq \subseteq D \times D$.

Recall that anti-symmetry requires $\forall x, y \in D: x \leq y \wedge y \leq x \Rightarrow x = y$.

• Binary relations can be understood as directed graphs:

$\{(a, a), (a, b), (b, c), (b, d), (d, c)\}$



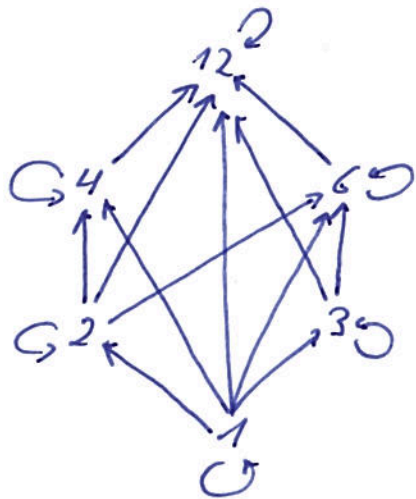
Partial orders yield particular graphs:

↳ Reflexivity = loops at every node

↳ Anti-symmetry = No cycles (except for loops)

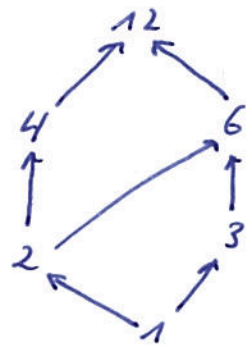
↳ Transitivity = Transitivity of the edges.

Example (Divisors of 12):



Hasse diagrams of partial orders do not draw

- loops and
- induced edges:



Definition (Join and meet):

Let (D, \leq) be a partial order.

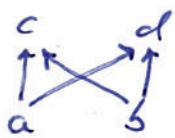
- The element $\sigma \in D$ is an upper bound of a set $X \subseteq D$, if $x \leq \sigma$ for all $x \in X$.

- The element $\sigma \in D$ is a least upper bound, also called join of $X \subseteq D$, if
 - σ is an upper bound of X and
 - $\sigma \leq \sigma'$ for every upper bound σ' of X .

We write $\sqcup X$ for the join of X .

- Similarly, $u = \sqcap X$ is the greatest lower bound, also called meet of X .

Example:



a and b:

- have c and d as upper bounds
- but do not have a least upper bound.

Definition (Complete lattice):

- A lattice is a partial order (D, \leq)

where join $a \sqcup b$ and meet $a \sqcap b$ exist for all $a, b \in D$.

- A lattice is complete if every subset $X \subseteq D$ has join and meet.

-> (*) The infix notation stands for $\sqcup\{a, b\}$.

Example: (1) $a \cdot b = \text{no lattice}$ (2) $\begin{matrix} \dots \\ 2 \\ 1 \\ 0 \end{matrix} \uparrow = \text{no complete lattice.}$

Lemma:

(1) Every complete lattice (D, \leq) has a unique

\hookrightarrow least element $\perp := \sqcup \emptyset = \prod D$

\hookrightarrow greatest element $\top := \prod \emptyset = \sqcup D$

(2) Every finite lattice (D, \leq) (with D finite) is complete.

16.2 Monotone Functions and Knaster and Tarski's Theorem:

Goal: Show the existence of fixed points
for monotone functions on complete lattices.

Definition (Monotone functions and fixed points):

Let (D, \leq) be a partial order.

• A function $f: D \rightarrow D$ is monotone,

\downarrow $x \leq y$ implies $f(x) \leq f(y)$ for all $x, y \in D$.

• A fixed point of $f: D \rightarrow D$ is an element $x \in D$
with $f(x) = x$.

We use $\text{Fix}(f) := \{x \in D \mid f(x) = x\}$ for the set of all fixed points of f .

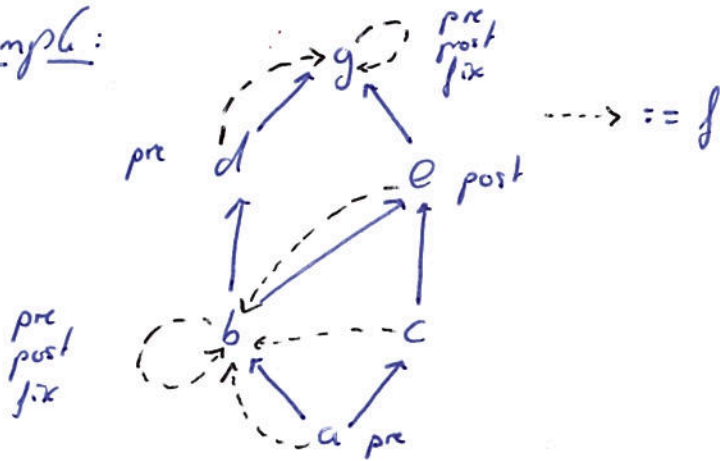
• A pre-fixed point is an element $x \in D$
with $x \leq f(x)$.

We use $\text{Prefix}(f)$ for the set of all pre-fixed points of f .

• A post-fixed point of f is an element $x \in D$
with $f(x) \leq x$.

We use $\text{Postfix}(f)$ for the set of all post-fixed points of f .

Example:



Theorem (Knaster and Tarski '55):

Let (D, \leq) be a complete lattice and let $f: D \rightarrow D$ be monotone.

(1) Then $\prod \text{Postfix}(f)$ is a fixed point of f .

Moreover, it is the least fixed point, and commonly denoted by $\text{lfp}(f)$.

(2) Similarly, $\text{l Prefix}(f)$ is the greatest fixed point of f , $\text{gfp}(f)$.

Proof:

We show (1), the reasoning for (2) is by duality (turn around the lattice).

• It is sufficient to show that $\prod \text{Postfix}(f)$ is a fixed point.

Since every fixed point is a post-fixed point,

$\prod \text{Postfix}(f)$ has to be smaller than every fixed point of f .

• We use $l := \prod \text{Postfix}(f)$ to denote the meet.

We first show

$$f(l) \leq l.$$

Since $l \leq l'$ for all $l' \in \text{Postfix}(f)$,

and since f is monotone, we obtain

$$f(l) \leq \underset{(\text{Postfix})}{f(l')} \leq l' \quad \text{for all } l' \in \text{Postfix}(f).$$

This means $f(l)$ is a lower bound for $\text{Postfix}(f)$.

Since l is the greatest lower bound for $\text{Postfix}(f)$,

$$f(l) \leq l$$

(*)

follows.

• We now show

$$l \leq f(l).$$

By (*) and monotonicity, we have

$$f(f(l)) \leq f(l).$$

But this means $f(l) \in \text{Postfix}(f)$.

Hence,

$$l \leq f(l).$$

(**)

• Anti-symmetry allows us to conclude $l = f(l)$ from (*) and (**). □

Note:

The above actually is a weak version of Knaster and Tarski's theorem. The full version states that $(\text{Fix}(f), \leq)$, i.e., the set of all fixed points, again forms a complete lattice.

16.3 Chains, Continuous Functions, and Kleene's Theorem

Goal: • Knaster and Tarski's result only tells us that fixed points exist (for monotone functions on complete lattices). • Kleene's theorem will allow us to compute them by iteration.

Definition:

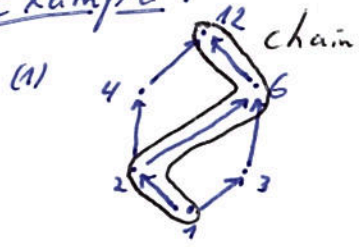
Let (D, \leq) be a partial order.

• A chain is a totally ordered subset $K \subseteq D$.

• We also give chains as sequences $(k_i)_{i \in \mathbb{N}}$ consisting of the members.

- A sequence $(k_i)_{i \in \mathbb{N}}$ is an ascending chain,
if $k_i \leq k_{i+1}$ for all $i \in \mathbb{N}$.
- A sequence $(k_i)_{i \in \mathbb{N}}$ is a descending chain,
if $k_i \geq k_{i+1}$ for all $i \in \mathbb{N}$.
- An ascending / descending chain $(k_i)_{i \in \mathbb{N}}$ becomes stationary,
if $\exists n \in \mathbb{N} \forall i \geq n: k_i = k_n$.
- One also says the chain stabilizes.
- (D, \leq) has finite height, if every chain $K \subseteq D$
has finitely many elements.
- (D, \leq) has bounded height, if there is an $n \in \mathbb{N}$
so that every chain $K \subseteq D$ has at most n elements.

Example:



(2) In (\mathbb{N}, \leq) , every descending chain becomes stationary.

Definition (Chain conditions):

- A partial order (D, \leq)
 - ↳ satisfies the ascending chain condition (ACC)
(one also says (D, \leq) is an Artinian lattice (Emil Artin, 1898-1962))
if every ascending chain $k_0 \leq k_1 \leq \dots$ becomes stationary;
 - ↳ satisfies the descending chain condition (DCC)
 (D, \leq) is Noetherian (Emmy Noether, 1882-1935))
if every descending chain $k_0 \geq k_1 \geq \dots$ becomes stationary.

Note: (ACC) and (DCC) are independent of the lattice conditions.

Lemma:

A partial order has finite height iff (ACC) and (DCC) hold.

Definition (Continuity):

Let (D, \leq) be a complete lattice.

A function $f: D \rightarrow D$ is

(i) \sqcup -continuous, \Leftrightarrow for every chain $K \subseteq D$:
(join-continuous, upward-continuous) $f(\sqcup K) = \sqcup f(K) := \sqcup \{f(k) \mid k \in K\}$.

(ii) \sqcap -continuous, \Leftrightarrow for every chain $K \subseteq D$:
(meet-continuous, downward-continuous) $f(\sqcap K) = \sqcap f(K) := \sqcap \{f(k) \mid k \in K\}$.

Our goal is to show that monotonicity implies \sqcup -continuity, provided the corresponding chain condition holds.

The proof relies on the existence of greatest and least elements.

Lemma:

Let (D, \leq) be a partial order.

- (i) If (ACC) holds, every non-empty chain $K \subseteq D$ has a greatest element $g \in K$ (with $k \leq g$ for all $k \in K$).
- (ii) If (DCC) holds, every non-empty chain $K \subseteq D$ contains a least element $l \in K$ (with $l \leq k$ for all $k \in K$).

Proof:

We show (i), (ii) is by duality.

Towards a contradiction, assume

$$\neg (\exists g \in K \forall k \in K: k \leq g)$$

which means

$$\forall g \in K \exists k \in K: g \text{ and } k \text{ are incomparable or } g < k. \quad (*)$$

Case (*) cannot hold:

Since both g and h stem from K , and since K is a chain, they have to be comparable.

We thus have

$$\forall g \in K \exists h \in K: g < h \quad (g \leq h \text{ and } g \neq h). \quad (**)$$

This contradicts (ACC), as can be seen from the following ascending chain that does not become stationary:

$k_0 :=$ an arbitrary element from K
 $k_{i+1} :=$ the $h \in K$ with $h > k_i$,
which exists by (**). □

Theorem (Monotonicity implies continuity):

Let (D, \leq) be a complete lattice and let $f: D \rightarrow D$ be monotone.

(i) If (D, \leq) satisfies (ACC), then f is \sqcup -continuous.

(ii) If (D, \leq) satisfies (DCC), then f is \sqcap -continuous.

Proof:

We again only show (i). Consider an ascending chain $K \in D$.

• To see that $\sqcup f(K) \leq f(\sqcup K)$,

note that $k \in \sqcup K$ for all $k \in K$.

Hence, $f(k) \leq f(\sqcup K)$ by monotonicity.

Since this holds for all $k \in K$,

$f(\sqcup K)$ is an upper bound for $f(K) = \{f(k) \mid k \in K\}$.

Hence, the least upper bound is smaller, $\sqcup f(K) \leq f(\sqcup K)$.

• To see that $f(\sqcup K) \leq \sqcup f(K)$,

note that $\sqcup K = g$ with g the greatest element from the above lemma

Now $f(\sqcup K) = f(g) \leq \sqcup f(K)$.

• The inequality holds as $g \in K$.

• We conclude by anti-symmetry. □

For Kleene's fixed point theorem, we consider chains obtained by iterating the function of interest on \perp . This gives us a way to actually compute the least fixed point.

Lemma:

Let (D, \leq) be a complete lattice and let $f: D \rightarrow D$ be monotone.

The sequence

$$(f^i(\perp))_{i \in \mathbb{N}} \quad \text{with } f^0(\perp) := \perp \text{ and } f^{i+1}(\perp) := f(f^i(\perp))$$

is an ascending chain.

Proof:

We proceed by induction and show that $f^i(\perp) \leq f^{i+2}(\perp)$ for all $i \in \mathbb{N}$.

Base case: $f^0(\perp) = \perp \leq f(\perp)$, since bottom is the least element of D , $\perp = \prod D$.

Induction step: Assume $f^i(\perp) \leq f^{i+2}(\perp)$.

Then

$$f^{i+1}(\perp) = f(f^i(\perp)) \leq \underset{\substack{\text{(IH +} \\ \text{Monotonicity)}}}{f(f^{i+2}(\perp))} = f^{i+2}(\perp).$$

□

Theorem (Kleene):

Let (D, \leq) be a complete lattice and $f: D \rightarrow D$ monotone.

(i) If f is \sqcup -continuous, $\text{lfp}(f) = \sqcup \{f^i(\perp) \mid i \in \mathbb{N}\}$.

(ii) If f is \prod -continuous, $\text{gfp}(f) = \prod \{f^i(\top) \mid i \in \mathbb{N}\}$.

Proof:

We again show (i).

The first step is to prove that $\sqcup \{f^i(\perp) \mid i \in \mathbb{N}\}$ is a fixed point:

$$f(\sqcup \{f^i(\perp) \mid i \in \mathbb{N}\}) \underset{\substack{\text{(}\sqcup\text{-continuity)}}}{=} \sqcup \{f^{i+1}(\perp) \mid i \in \mathbb{N}\} \underset{\substack{\text{(}\perp = \prod D)}}{=} \sqcup \{f^i(\perp) \mid i \in \mathbb{N}\}.$$

• We now show that $\bigcup \{f^i(\perp) \mid i \in \mathbb{N}\}$ is the least fixed point.
 To this end, we consider $d \in D$ with $f(d) = d$
 and show that $\bigcup \{f^i(\perp) \mid i \in \mathbb{N}\}$ is smaller.

• By an induction on $i \in \mathbb{N}$, we prove $f^i(\perp) \leq d$ for all $i \in \mathbb{N}$.

Base case: $f^0(\perp) = \perp \leq d$, since $\perp = \prod D$.

Induction step: Assume $f^i(\perp) \leq d$.

We then get

$$f^{i+1}(\perp) = f(f^i(\perp)) \leq f(d) = d.$$

(IH +
Monotonicity)

• Since $f^i(\perp) \leq d$ for all $i \in \mathbb{N}$, d is an upper bound for $\{f^i(\perp) \mid i \in \mathbb{N}\}$,
 and thus $\bigcup \{f^i(\perp) \mid i \in \mathbb{N}\} \leq d$. □

In combination with the above theorem on monotonicity implying continuity under appropriate chain conditions, we obtain the following.

Corollary:

Let (D, \leq) be a complete lattice with (FCC) and (DCC).

Let $f: D \rightarrow D$ be monotone.

Then • $\text{lfp}(f) = \bigcup \{f^i(\perp) \mid i \in \mathbb{N}\}$
 $= f^\omega(\perp)$ with $f^\omega(\perp) = f^{\omega+1}(\perp)$.

• $\text{gfp}(f) = \bigcap \{f^i(\top) \mid i \in \mathbb{N}\}$
 $= f^\omega(\top)$ with $f^\omega(\top) = f^{\omega+1}(\top)$.

Recall that (FCC) and (DCC) hold iff the lattice has finite height.